

#### Tuning the Storage Subsystem

## The Controller Cache

- Read Cache: performs read-ahead
  - after read request, controller continues to read and store in cache
  - database can do better read-ahead since it knows the access patterns
  - in general it is better to turn read cache off!
- Write-back mode: request terminates when data is written to cache
  - data is written from cache to disk later
  - writes become faster since they do not have to wait for the disk
  - if cache contents get lost (power failure, no battery), then data is lost
- Write-through mode: request terminates when data is written to disk
  - if cache has not battery, this mode is safer

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• if cache is overloaded, write-through might be more efficient (depends on the efficiency of the replacement policy algorithm)

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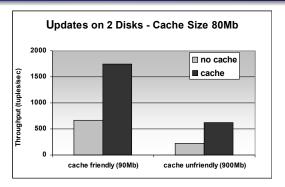
RAID – Redundant Arrays of Inexpensive Disks

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## • RAID array

- multiple hard disks
- one logical disk (operating system sees only one disk)
- disks divide and replicate data
- RAID controller is the interface
- Benefits:
  - fault tolerance by introducing redundancy
  - increased throughput due to parallel disk access

# Tuning the Storage Subsystem Write-Back Mode – Experiment



- Cache controller in write-back mode vs. no cache.
- Cache friendly load: volume of update slightly larger than cache.
- Cache unfriendly: volume of update much larger than cache.
- Write-back gives similar benefit in both cases.
- The controller implements an efficient replacement policy.

#### SQL Server 7 on Windows 2000 DBT – Hardware Tuning

#### Tuning the Storage Subsystem

#### RAID Levels

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- RAID configurations are numbered ("levels"):
  - RAID 0: striping
  - RAID 1: mirroring
  - RAID 5: rotated parity striping
  - RAID 10: striped mirroring
- other (less important) RAID levels exist

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#### Tuning the Storage Subsystem RAID 0 - Striping

#### • RAID 0 - striping:

- data is split into stripes of the same size
- consecutive stripes are written onto consecutive disks
- Example: stripe size 1kB, 2 disks
  - write datablock A of 8kB
  - block is split into  $A = A_1 + A_2 + \ldots + A_8$
  - disk 1:  $A_1, A_3, A_5, A_7$ , disk 2:  $A_2, A_4, A_6, A_8$
- Read/Write: RAID 0 with n disks, stripe size s
  - small data  $\leq$  s: read/write results in access to one physical disk
  - large data  $\geq s \times n$ : read/write results in parallel access to n disks
- Benefits/drawbacks:

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- + fast sequential read/write and concurrent seeks
- + 100% utilization of disk space (=cheap)

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- no fault-tolerance (array inaccessible if single disk fails)
- Database use: temporary files

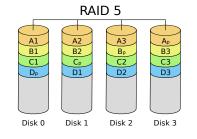
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## RAID 5 – Rotated Parity Striping

- RAID 5 rotated parity striping: n disks
  - fault tolerance by error correction (instead of full redundancy)

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- striped as in RAID 0, but n-1 stripes have an additional parity stripe
- parity stripes are evenly distributed over disks
- Example: stripe size 1kB, 4 disks
  - write datablock *AB* of 6*kB*
  - datablock is split into  $AB = A_1 + A_2 + A_3 + A_p + B_1 + B_2 + B_3 + B_p$
  - disk 1: A<sub>1</sub>, B<sub>1</sub>, disk 2: A<sub>2</sub>, B<sub>2</sub>, disk 3: A<sub>3</sub>, B<sub>p</sub>, disk 4: A<sub>p</sub>, B<sub>3</sub>



## RAID 1 – Mirroring

- RAID 1 mirroring: 2 disks
  - the same data is written to both disks ("mirrored")
  - no striping
- Example: 2 disks
  - write datablocks  $A_1, A_2, A_3, A_4$
  - disk 1: writes  $A_{1-4}$ , disk 2: writes  $A_{1-4}$

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- Write one data block:
  - physical write two both disks
  - operation terminates when slower disk is done (!)
- Read of one data block: physical read from single (least busy) disks

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• Benefits/drawbacks:

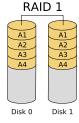
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- + fault tolerant (no interruption if one disk fails)
- + concurrent seeks (faster random read access)
- only 50% utilization of disk space (=expensive)
- write speed not increased from single-disk solution
- Database use: log file (fault tolerance, sequential writes)

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### RAID 5 – Rotated Parity Striping

- Read: like RAID 0 with n-1 disks, parity stripe is not read
- Writing 1 stripe requires 2 physical reads and writes
- Write: Update data stripe S with S'
  - 1. read old data stripe S and old parity stripe P
  - 2. new parity stripe  $P' = S \operatorname{xor} S' \operatorname{xor} P$ 
    - (for each bit flipped between S and S' flip the corresponding bit in P)
  - 3. write S' and P' (substituting S and P, respectively)
- Recovery: *n* disks, failure on disk *x* 
  - S<sub>i</sub> is the stripe on disk i (either parity or data)
  - lost stripe  $S_x = S_1 \operatorname{xor} \ldots \operatorname{xor} S_{x-1} \operatorname{xor} S_{x+1} \operatorname{xor} \ldots \operatorname{xor} S_n$
- Benefits/drawbacks:
  - $+\,$  fault tolerant (slowdown, but no interruption if one disk fails)
  - + fast sequential read and concurrent seeks
  - + (100 100/n)% utilization of disk space (=cheap)
  - write is slower than single-disk solution
  - recovery after failure much more difficult than with RAID 1  $\,$
- Database use: data and index files (if reads predominate writes)



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RAID 0

A1 A3 A5 A7 A8

Disk 0

Disk 1

## RAID 10 – Mirroring + Striping

- RAID 1+0: mirrored RAID 0
  - stripe data on first n/2 disks (as in RAID 0)

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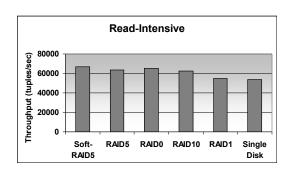
• use the other disks to mirror these disks

#### • Benefits/drawbacks:

- + best performance of all RAID levels
- + fault tolerant (no interruption if one disk fails)
- + fast sequential read/write and concurrent seeks
- 50% utilization of disk space (=expensive)
- Database use:
  - log file if RAID 1 is too slow
  - data and index files if writes are too slow on RAID 5

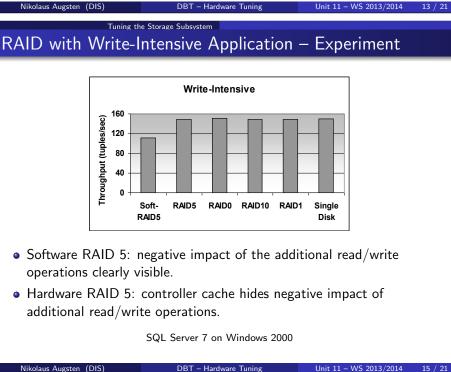
## RAID with Read-Intensive Application – Experiment

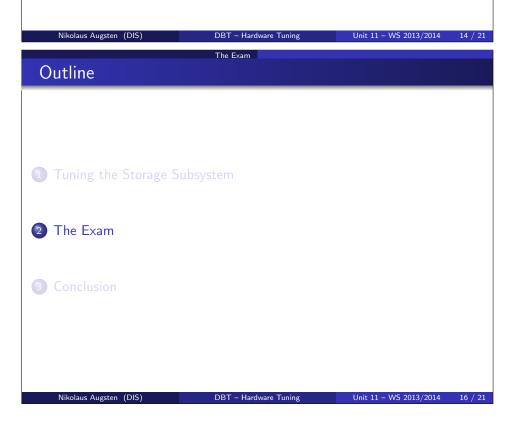
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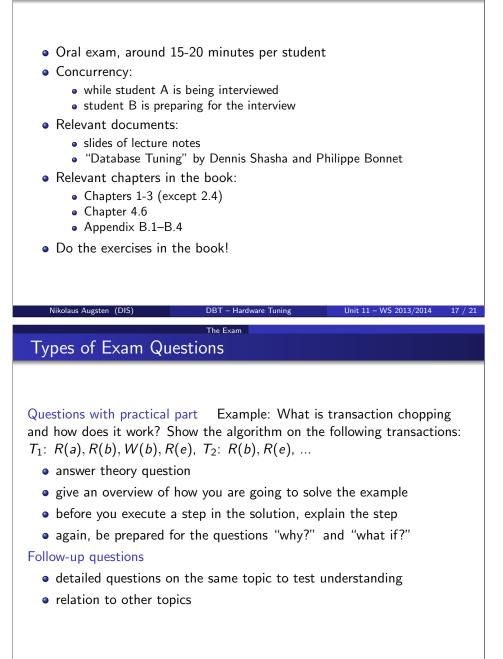
- RAID 1 slightly improves on single disk solution.
- Striped RAID levels (0,5,10) significantly improve read performance.

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#### Exam



The Exam

## Types of Exam Questions

Theory question Example: Explain write-ahead logging and how the logging mechanism can be tuned.

The Evam

- illustrate situation (database buffer, log buffer, log file, data files)
- use correct terminology and give precise definitions (e.g., what is a checkpoint?)

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Conclusion

- structure your answer (how does WAL work? list tuning opportunities, then discuss each of them)
- discussion (advantage/disadvantage)
- be prepared for the questions "why?" and "what if?"

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Outline

Tuning the Storage Subsystem

2 The Exam

3 Conclusion

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# Summary

- Tuning the Storage Subsystem
  - raw vs. cooked file
  - setting the controller cache mode

Conclusion

• choosing the RAID level

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