

# IT Security

## Multilevel Databases

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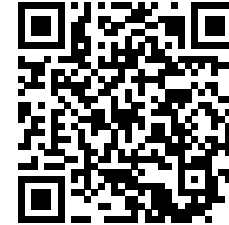
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All infos about the database part in this lecture

<http://dbresearch.uni-salzburg.at/teaching/2015ws/its/>



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MAC: Mandatory Access Control

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## Mandatory Access Control (MAC)

- Why is discretionary access control (DAC) not enough?
  - users have the freedom to give other users access to data
  - all users see the same data (if they have access)
  - security policies cannot be centrally enforced
- Some applications need multilevel security
  - government, military, intelligence service
  - many industrial and corporate applications
- MAC is implemented in some DBMS (e.g., Oracle Label Security since 2009) or special versions of DBMS (e.g., SE-PostgreSQL)
- also operating systems implement MAC (SE-Linux, Windows Vista and later)

## MAC Basics

- Security classes: levels of trust
  - TS (top secret) > S (secret) > C (confidential) > U (unclassified, public)
- Subjects  $s$ 
  - users, roles, accounts, programs
  - clearance  $clear(s)$  is the trustworthiness of  $s$
  - $clear(s)$  is a security class
- Objects  $o$ :
  - data objects (e.g., relation, tuple, attribute values)
  - classification  $class(o)$  is the sensitivity of the data object
  - $class(o)$  is a security class

## Bell LaPadula

- Example of MAC used in database (and many other) systems
- Named after developers D. E. Bell and L. J. LaPadula
- Access control rules
  - no read-up:  $s$  is allowed to read  $o$  only if  $clear(s) \geq class(o)$
  - no write-down:  $s$  is allowed to write  $o$  only if  $clear(s) \leq class(o)$  (also called \*-property)
  - respect DAC: respect discretionary access control rules
- Trusted subjects
  - must be trustworthy according to security policy
  - not restricted by the \*-property
  - can transfer data from higher to lower sensitivity

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## Multilevel Model

- **Multilevel relation**

- each attribute and each tuple in  $R(A_1, A_2, \dots, A_n)$  are classified
- $C_i = \text{class}(A_i)$  is an attribute classification
- $TC \geq \max\{C_i \mid 1 \leq i \leq n\}$  is the tuple classification
- the schema of the multilevel relation is

$$R(A_1, C_1, A_2, C_2, \dots, A_n, C_n, TC)$$

## Reading from Multilevel Relations

- **Security requirement**

- users should not even know which data they cannot access
- system should not reject requests for non-authorized data
- but still the user should see a consistent view of the table

- Each clearance class  $c$  sees a different **instance**  $R^c$  of  $R$ :

$$R^c = (A_1^c, C_1^c, A_2^c, C_2^c, \dots, A_n^c, C_n^c, TC^c)$$

- **Attributes**  $A_i^c$  visible by  $s$  with  $\text{clear}(s) = c$ :

- $A_i^c = A_i$  if  $C_i \leq c$
- $A_i^c = \text{NULL}$  if  $C_i > c$

- **Classifications**  $C_i^c$  and  $TC^c$ :

- $C_i^c = \min\{C_i, c\}$
- $TC^c = \min\{TC, c\}$

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## How to Deal with Updates?

- **Problem:**

- subject with low clearance sees NULL value and tries to change it
- but this NULL value is due to the low clearance

- **Option 1 (bad):** update value

- values of subjects with higher clearance get lost
- writers do not even realize that they are doing something harmful

- **Option 2 (bad):** reject update

- writing subject can infer that there is a sensitive non-NULL value
- can be systematically exploited

- **Option 3 (good):** Polyinstantiation

- maintain multiple versions of tuples
- versioned tuples must differ by sensitivity class  $TC$
- new model for integrity is required!

## Integrity in Multi-Level Databases

- Entity integrity
- Null integrity
- Inter-instance integrity
- Polyinstantiation integrity

## Entity Integrity

- Keys in instance  $R^c$  are called **apparent key**
- **Entity integrity**: for each  $R^c$  and for each tuple in  $R^c$ 
  1. key values must not be NULL
  2. all key attributes must have identical sensitivity class
  3. non-keys must be at least as sensitive as key

## Null Integrity

- **Null integrity**: for each  $R^c$ 
  1. NULL values always have sensitivity of key
  2. freedom of subsumption (= no unnecessary tuples)

## Inter Instance Integrity

- **Inter instance integrity**: for any pair  $R^c, R^{c'}$  with  $c' < c$

$$R^{c'} = f(R^c)$$

where  $f$  is called **filter**.

- The filter has the following properties
  1. for each tuple in  $R^c$  with key visible by  $c'$  a tuple must exist in  $R^{c'}$
  2. no other tuples exist in  $R^{c'}$
  3. subsumed tuples are eliminated

## Polyinstantiation Integrity

- Polyinstantiation integrity: uniqueness of tuples in  $R$ 
  - functional dependency

$$(key, C_{key}, C_i) \rightarrow A_i$$

must hold for any  $A_i$  in instance  $R^c$

## Implementation of Multilevel Databases

- integrity constraints allow implementation on top of “normal” relational system
- multilevel relation is fragmented into normal relations
- user queries compute answer from fragmented relations

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## Example: Oracle Label Security/1

- Label security provides MAC for Oracle DBMS
- Each tuple and each user has a label
- Labels consist of
  - level (class / clearance)
  - compartments: segregate data within a given level
  - groups: segregate data within level using organizational hierarchy
- subject  $s$  can access object  $o$  if
  - label of  $s$  must be at least level of  $o$
  - $s$  must have all compartments of  $o$
  - $s$  must have at least one group or supergroup of  $o$

## User labels

- User labels
  - max read clearance
  - min write clearance
  - default clearance (at login)
  - row level: default for inserted tuples
  - read and write compartments
  - read and write groups
- Trusted users / stored procedures
  - read / writeup / writedown
  - write across: change compartment and group
  - profile access: become other user (like Unix 'su')