

# Database Tuning

## Hardware Tuning

Nikolaus Augsten

nikolaus.augsten@sbg.ac.at  
Department of Computer Sciences  
University of Salzburg



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Adapted from “Database Tuning” by Dennis Shasha and Philippe Bonnet.

# Outline

- 1 Tuning the Storage Subsystem
- 2 The Exam
- 3 Conclusion

# Overview

- Tuning the storage subsystem involves configuring:
  - disk allocation
  - disk array (RAID level)
  - controller cache

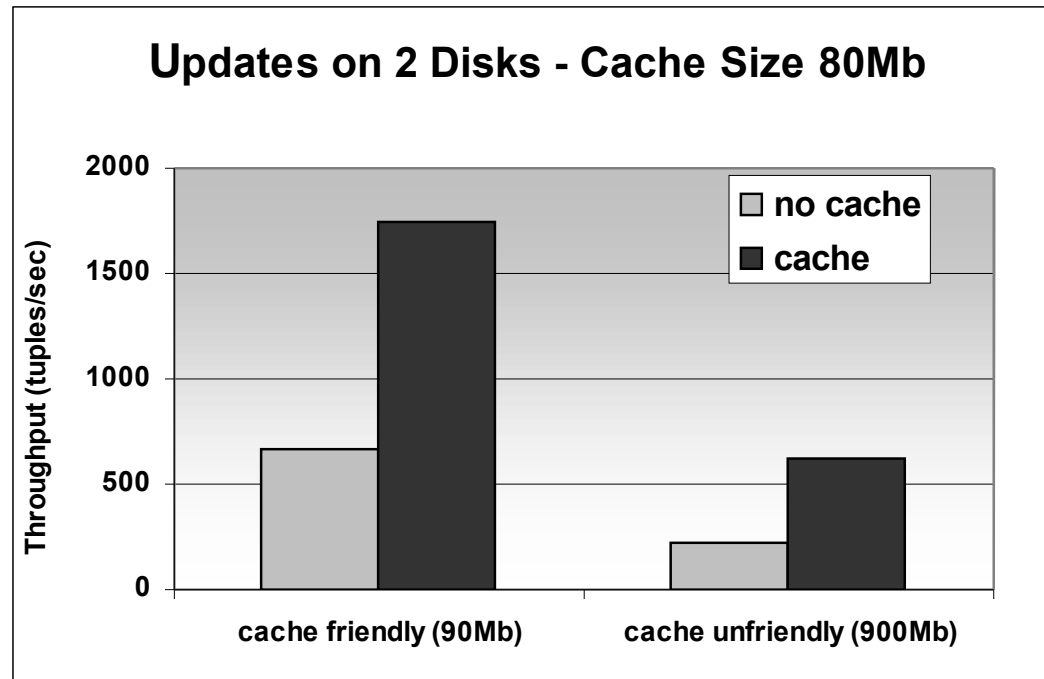
# Disk Allocation – Raw vs. Cooked Files

- Cooked file:
  - regular operating system file
  - buffered through the operating system
  - logically contiguous blocks might not be physically contiguous
  - possibly indirection to access a particular page (inodes in Unix/Linux)
- Raw file: also “character special device”
  - block device (hard disk) configured as raw device (using raw command)  
e.g., `/dev/rhd0f` is the raw device of block device `/dev/hd0f`
  - not buffered by the operating system
  - logically contiguous blocks are physically contiguous
  - more efficient than cooked files

# The Controller Cache

- **Read Cache:** performs read-ahead
  - after read request, controller continues to read and store in cache
  - database can do better read-ahead since it knows the access patterns
  - in general it is better to turn read cache off!
- **Write-back mode:** request terminates when data is written to cache
  - data is written from cache to disk later
  - writes become faster since they do not have to wait for the disk
  - if cache contents get lost (power failure, no battery), then data is lost
- **Write-through mode:** request terminates when data is written to disk
  - if cache has no battery, this mode is safer
  - if cache is overloaded, write-through might be more efficient (depends on the efficiency of the replacement policy algorithm)

# Write-Back Mode – Experiment



- Cache controller in write-back mode vs. no cache.
- Cache friendly load: volume of update slightly larger than cache.
- Cache unfriendly: volume of update much larger than cache.
- Write-back gives similar benefit in both cases.
- The controller implements an efficient replacement policy.

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# RAID – Redundant Arrays of Inexpensive Disks

- RAID array
  - multiple hard disks
  - one logical disk (operating system sees only one disk)
  - disks divide and replicate data
  - RAID controller is the interface
- Benefits:
  - fault tolerance by introducing redundancy
  - increased throughput due to parallel disk access

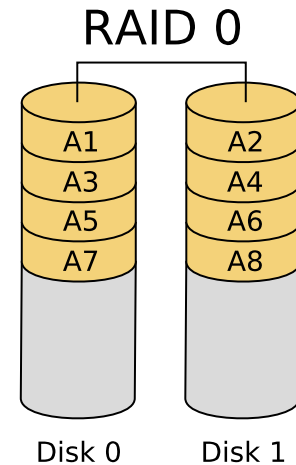
# RAID Levels

- RAID configurations are numbered (“levels”):
  - RAID 0: striping
  - RAID 1: mirroring
  - RAID 5: rotated parity striping
  - RAID 10: striped mirroring
- other (less important) RAID levels exist



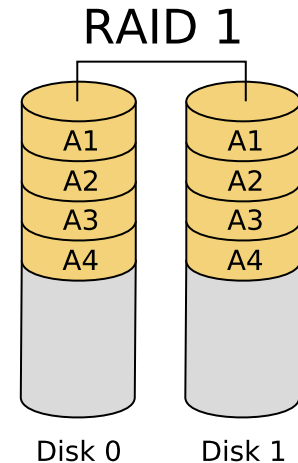
# RAID 0 – Striping

- RAID 0 – striping:
  - data is split into **stripes** of the same size
  - consecutive stripes are written onto consecutive disks
- Example: stripe size 1kB, 2 disks
  - write datablock  $A$  of 8kB
  - block is split into  $A = A_1 + A_2 + \dots + A_8$
  - disk 1:  $A_1, A_3, A_5, A_7$ , disk 2:  $A_2, A_4, A_6, A_8$
- Read/Write: RAID 0 with  $n$  disks, stripe size  $s$ 
  - small data  $\leq s$ : read/write results in access to one physical disk
  - large data  $\geq s \times n$ : read/write results in parallel access to  $n$  disks
- Benefits/drawbacks:
  - + fast sequential read/write and concurrent seeks
  - + 100% utilization of disk space (=cheap)
  - no fault-tolerance (array inaccessible if single disk fails)
- Database use: temporary files



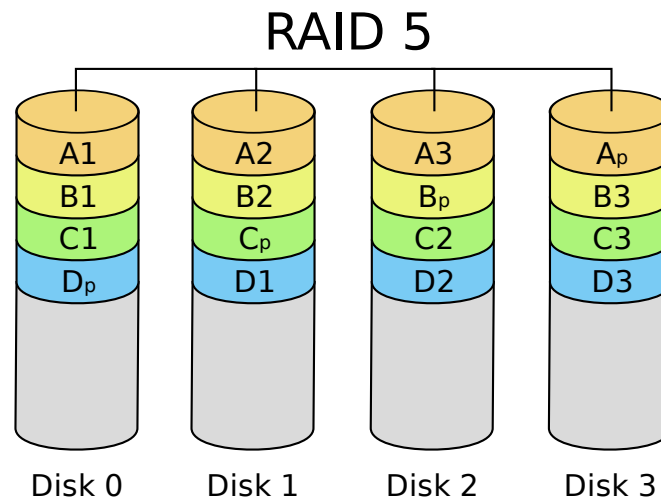
# RAID 1 – Mirroring

- RAID 1 – mirroring: 2 disks
  - the same data is written to both disks (“mirrored”)
  - no striping
- Example: 2 disks
  - write datablocks  $A_1, A_2, A_3, A_4$
  - disk 1: writes  $A_{1-4}$ , disk 2: writes  $A_{1-4}$
- Write one data block:
  - physical write two both disks
  - operation terminates when slower disk is done (!)
- Read of one data block: physical read from single (least busy) disks
- Benefits/drawbacks:
  - + fault tolerant (no interruption if one disk fails)
  - + concurrent seeks (faster random read access)
  - only 50% utilization of disk space (=expensive)
  - write speed not increased from single-disk solution
- Database use: log file (fault tolerance, sequential writes)



# RAID 5 – Rotated Parity Striping

- RAID 5 – rotated parity striping:  $n$  disks
  - fault tolerance by error correction (instead of full redundancy)
  - striped as in RAID 0, but  $n - 1$  stripes have an additional parity stripe
  - parity stripes are evenly distributed over disks
- Example: stripe size 1kB, 4 disks
  - write datablock  $AB$  of 6kB
  - datablock is split into  $AB = A_1 + A_2 + A_3 + A_p + B_1 + B_2 + B_3 + B_p$
  - disk 1:  $A_1, B_1$ , disk 2:  $A_2, B_2$ , disk 3:  $A_3, B_p$ , disk 4:  $A_p, B_3$



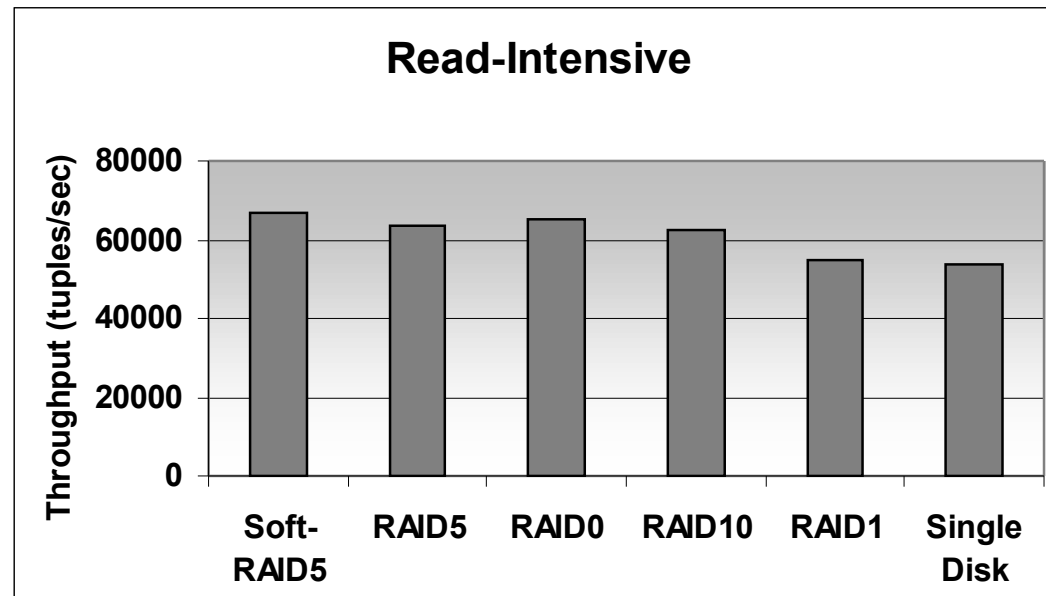
# RAID 5 – Rotated Parity Striping

- Read: like RAID 0 with  $n - 1$  disks, parity stripe is not read
- Writing 1 stripe requires 2 physical reads and writes
- Write: Update data stripe  $S$  with  $S'$ 
  1. read old data stripe  $S$  and old parity stripe  $P$
  2. new parity stripe  $P' = S \text{ xor } S' \text{ xor } P$   
(for each bit flipped between  $S$  and  $S'$  flip the corresponding bit in  $P$ )
  3. write  $S'$  and  $P'$  (substituting  $S$  and  $P$ , respectively)
- Recovery:  $n$  disks, failure on disk  $x$ 
  - $S_i$  is the stripe on disk  $i$  (either parity or data)
  - lost stripe  $S_x = S_1 \text{ xor } \dots \text{ xor } S_{x-1} \text{ xor } S_{x+1} \text{ xor } \dots \text{ xor } S_n$
- Benefits/drawbacks:
  - + fault tolerant (slowdown, but no interruption if one disk fails)
  - + fast sequential read and concurrent seeks
  - +  $(100 - 100/n)\%$  utilization of disk space (=cheap)
    - write is slower than single-disk solution
    - recovery after failure much more difficult than with RAID 1
- Database use: data and index files (if reads predominate writes)

# RAID 10 – Mirroring + Striping

- RAID 1+0: mirrored RAID 0
  - stripe data on first  $n/2$  disks (as in RAID 0)
  - use the other disks to mirror these disks
- Benefits/drawbacks:
  - + best performance of all RAID levels
  - + fault tolerant (no interruption if one disk fails)
  - + fast sequential read/write and concurrent seeks
    - 50% utilization of disk space (=expensive)
- Database use:
  - log file if RAID 1 is too slow
  - data and index files if writes are too slow on RAID 5

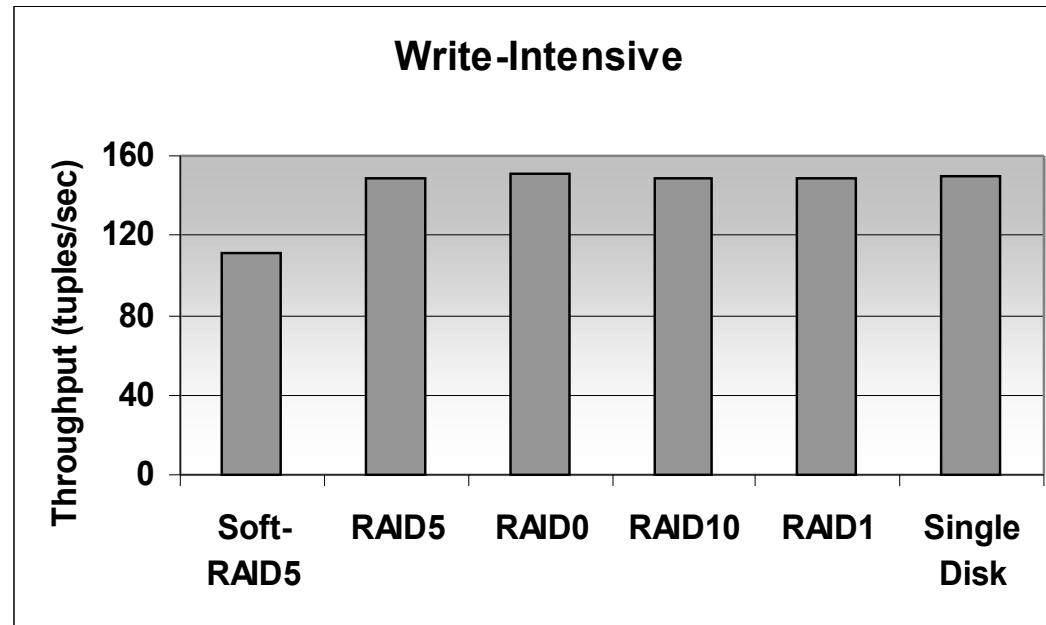
# RAID with Read-Intensive Application – Experiment



- RAID 1 slightly improves on single disk solution.
- Striped RAID levels (0,5,10) significantly improve read performance.

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# RAID with Write-Intensive Application – Experiment



- Software RAID 5: negative impact of the additional read/write operations clearly visible.
- Hardware RAID 5: controller cache hides negative impact of additional read/write operations.

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# Exam

- Oral exam, around 15-20 minutes per student
- Concurrency:
  - while student A is being interviewed
  - student B is preparing for the interview
- Relevant documents:
  - slides of lecture notes
  - “Database Tuning” by Dennis Shasha and Philippe Bonnet
- Relevant chapters in the book:
  - Chapters 1-3 (except 2.4)
  - Chapter 4.6
  - Appendix B.1–B.4
- Do the exercises in the book!

# Types of Exam Questions

**Theory question** Example: Explain write-ahead logging and how the logging mechanism can be tuned.

- illustrate situation (database buffer, log buffer, log file, data files)
- use correct terminology and give precise definitions (e.g., what is a checkpoint?)
- structure your answer (how does WAL work? list tuning opportunities, then discuss each of them)
- discussion (advantage/disadvantage)
- be prepared for the questions “why?” and “what if?”

# Types of Exam Questions

**Questions with practical part** Example: What is transaction chopping and how does it work? Show the algorithm on the following transactions:  $T_1$ :  $R(a), R(b), W(b), R(e)$ ,  $T_2$ :  $R(b), R(e), \dots$

- answer theory question
- give an overview of how you are going to solve the example
- before you execute a step in the solution, explain the step
- again, be prepared for the questions “why?” and “what if?”

## Follow-up questions

- detailed questions on the same topic to test understanding
- relation to other topics
- possibly additional questions on other topics

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# Summary

- Tuning the Storage Subsystem
  - raw vs. cooked file
  - setting the controller cache mode
  - choosing the RAID level