Database Tuning Hardware Tuning

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Adapted from "Database Tuning" by Dennis Shasha and Philippe Bonnet.

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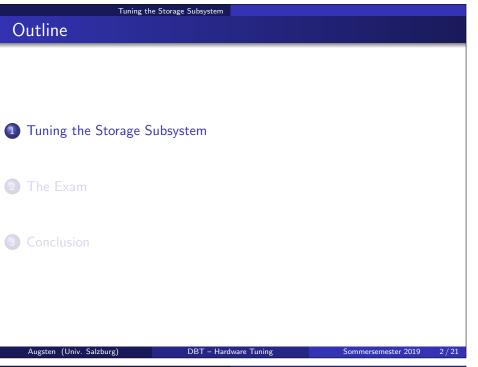
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Tuning the Storage Subsystem

Overview

- Tuning the storage subsystem involves configuring:
 - disk allocation
 - disk array (RAID level)
 - controller cache



Tuning the Storage Subsystem

Disk Allocation - Raw vs. Cooked Files

- Cooked file:
 - regular operating system file
 - buffered through the operating system
 - logically contiguous blocks might not be physically contiguous
 - possibly indirection to access a particular page (inodes in Unix/Linux)
- Raw file: also "character special device"
 - block device (hard disk) configured as raw device (using raw command) e.g., /dev/rsd0f is the raw device of block device /dev/sd0f
 - not buffered by the operating system
 - logically contiguous blocks are physically contiguous
 - more efficient than cooked files

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The Controller Cache

- Read Cache: performs read-ahead
 - after read request, controller continues to read and store in cache
 - database can do better read-ahead since it knows the access patterns
 - in general it is better to turn read cache off!
- Write-back mode: request terminates when data is written to cache
 - data is written from cache to disk later
 - writes become faster since they do not have to wait for the disk
 - if cache contents get lost (power failure, no battery), then data is lost
- Write-through mode: request terminates when data is written to disk
 - if cache has no battery, this mode is safer
 - if cache is overloaded, write-through might be more efficient (depends on the efficiency of the replacement policy algorithm)

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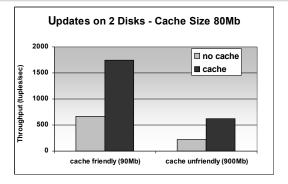
Tuning the Storage Subsystem

RAID - Redundant Arrays of Inexpensive Disks

- RAID array
 - multiple hard disks
 - one logical disk (operating system sees only one disk)
 - disks divide and replicate data
 - RAID controller is the interface.
- Benefits:
 - fault tolerance by introducing redundancy
 - increased throughput due to parallel disk access

Tuning the Storage Subsystem

Write-Back Mode – Experiment



- Cache controller in write-back mode vs. no cache.
- Cache friendly load: volume of update slightly larger than cache.
- Cache unfriendly: volume of update much larger than cache.
- Write-back gives similar benefit in both cases.
- The controller implements an efficient replacement policy.

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Tuning the Storage Subsystem **RAID** Levels

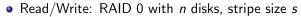
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- RAID configurations are numbered ("levels"):
 - RAID 0: striping • RAID 1: mirroring
 - RAID 5: rotated parity striping
 - RAID 10: striped mirroring
- other (less important) RAID levels exist

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RAID 0 - Striping

- RAID 0 striping:
 - data is split into stripes of the same size
 - consecutive stripes are written onto consecutive disks
- Example: stripe size 1kB, 2 disks
 - write datablock A of 8kB
 - block is split into $A = A_1 + A_2 + \ldots + A_8$
 - disk 1: A_1, A_3, A_5, A_7 , disk 2: A_2, A_4, A_6, A_8



- small data < s: read/write results in access to one physical disk
- large data $\geq s \times n$: read/write results in parallel access to n disks
- Benefits/drawbacks:
 - + fast sequential read/write and concurrent seeks
 - + 100% utilization of disk space (=cheap)
 - no fault-tolerance (array inaccessible if single disk fails)
- Database use: temporary files

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RAID 0

A1 A3 A5 A7

Disk 0

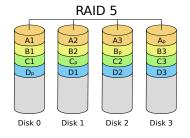
A2 A4 A6 A8

Disk 1

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RAID 5 - Rotated Parity Striping

- RAID 5 rotated parity striping: *n* disks
 - fault tolerance by error correction (instead of full redundancy)
 - striped as in RAID 0, but n-1 stripes have an additional parity stripe
 - parity stripes are evenly distributed over disks
- Example: stripe size 1kB, 4 disks
 - write datablock AB of 6kB
 - datablock is split into $AB = A_1 + A_2 + A_3 + A_p + B_1 + B_2 + B_3 + B_p$
 - disk 1: A_1 , B_1 , disk 2: A_2 , B_2 , disk 3: A_3 , B_p , disk 4: A_p , B_3



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RAID 1 - Mirroring

- RAID 1 mirroring: 2 disks
 - the same data is written to both disks ("mirrored")
 - no striping
- Example: 2 disks
 - write datablocks A_1 , A_2 , A_3 , A_4
 - disk 1: writes A_{1-4} , disk 2: writes A_{1-4}
- Write one data block:
 - physical write two both disks
 - operation terminates when slower disk is done (!)
- Read of one data block: physical read from single (least busy) disks
- Benefits/drawbacks:
 - + fault tolerant (no interruption if one disk fails)
 - + concurrent seeks (faster random read access)
 - only 50% utilization of disk space (=expensive)
 - write speed not increased from single-disk solution
- Database use: log file (fault tolerance, sequential writes)

RAID 1

Disk 0

A1 A2 A3 A4

Disk 1

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RAID 5 - Rotated Parity Striping

- Read: like RAID 0 with n-1 disks, parity stripe is not read
- Writing 1 stripe requires 2 physical reads and writes
- Write: Update data stripe S with S'
 - read old data stripe S and old parity stripe P
 - a new parity stripe $P' = S \times S' \times S' \times P$ (for each bit flipped between S and S' flip the corresponding bit in P)
 - write S' and P' (substituting S and P, respectively)
- Recovery: n disks, failure on disk x
 - S_i is the stripe on disk i (either parity or data)
 - lost stripe $S_x = S_1 \operatorname{xor} \ldots \operatorname{xor} S_{x-1} \operatorname{xor} S_{x+1} \operatorname{xor} \ldots \operatorname{xor} S_n$
- Benefits/drawbacks:
 - + fault tolerant (slowdown, but no interruption if one disk fails)
 - + fast sequential read and concurrent seeks
 - + (100 100/n)% utilization of disk space (=cheap)
 - write is slower than single-disk solution
 - recovery after failure much more difficult than with RAID 1
- Database use: data and index files (if reads predominate writes)

Tuning the Storage Subsystem

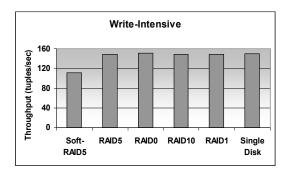
RAID 10 - Mirroring + Striping

- RAID 1+0: mirrored RAID 0
 - stripe data on first n/2 disks (as in RAID 0)
 - use the other disks to mirror these disks
- Benefits/drawbacks:
 - + best performance of all RAID levels
 - + fault tolerant (no interruption if one disk fails)
 - + fast sequential read/write and concurrent seeks
 - 50% utilization of disk space (=expensive)
- Database use:
 - log file if RAID 1 is too slow
 - data and index files if writes are too slow on RAID 5

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Tuning the Storage Subsystem RAID with Write-Intensive Application – Experiment

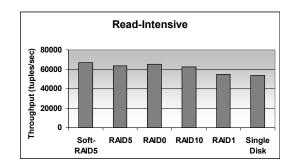


- Software RAID 5: negative impact of the additional read/write operations clearly visible.
- Hardware RAID 5: controller cache hides negative impact of additional read/write operations.

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Tuning the Storage Subsystem

RAID with Read-Intensive Application – Experiment



- RAID 1 slightly improves on single disk solution.
- Striped RAID levels (0,5,10) significantly improve read performance.

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The Exam

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Outline

Tuning the Storage Subsystem

2 The Exam

Conclusion

The Exam

Exam

- Oral exam, around 15-20 minutes per student
- Concurrency:
 - while student A is being interviewed
 - student B is preparing for the interview
- Relevant documents:
 - slides of lecture notes
 - "Database Tuning" by Dennis Shasha and Philippe Bonnet
- Relevant chapters in the book:
 - Chapters 1-3 (except 2.4)
 - Chapter 4.6
 - Appendix B.1–B.4
- Do the exercises in the book!

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The Exam

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Types of Exam Questions

Questions with practical part Example: What is transaction chopping and how does it work? Show the algorithm on the following transactions: T_1 : $R(a), R(b), W(b), R(e), T_2: R(b), R(e), ...$

- answer theory question
- give an overview of how you are going to solve the example
- before you execute a step in the solution, explain the step
- again, be prepared for the questions "why?" and "what if?"

Follow-up questions

- detailed questions on the same topic to test understanding
- relation to other topics
- possibly additional questions on other topics

Types of Exam Questions

Theory question Example: Explain write-ahead logging and how the logging mechanism can be tuned.

- illustrate situation (database buffer, log buffer, log file, data files)
- use correct terminology and give precise definitions (e.g., what is a checkpoint?)
- structure your answer (how does WAL work? list tuning opportunities, then discuss each of them)
- discussion (advantage/disadvantage)
- be prepared for the questions "why?" and "what if?"

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Conclusion

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Outline

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Summary

• Tuning the Storage Subsystem

• raw vs. cooked file

• setting the controller cache mode

• choosing the RAID level

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