

DISTRIBUTED LOCKING PROTOCOLS – OVERVIEW

| | Single-Lock Manager | Primary Copy | Majority | Biased |
|---------------------------|---------------------------------------|--|--|--|
| Deadlock Handling | Centralized | Distributed | Distributed | Distributed |
| Availability (*) | global SPoF (central lock manager) | SPoF per data item Q (primary copy of Q) | No SPoF $(\lfloor \frac{n}{2} \rfloor + 1$ replicas must be up) | Read: No SPoF Write: all n replicas of Q must be up |
| Bottleneck | Yes | No | No | No |
| Supports Replicas | Yes | Yes | Yes | Yes |
| # of Messages | Read: 2 Write: 2 Unlock: 1 | Read: 2 Write: 2 Unlock: 1 | Read/Write: $2 (\lfloor \frac{n}{2} \rfloor + 1)$ Unlock: $(\lfloor \frac{n}{2} \rfloor + 1)$ | Read: 2 Write: $2n$ Unlock: Read: 1, Write: n |
| Read from Replicas | Any | Any | Any | Any |
| Write to Replicas | All | All | All | All |

SPoF Single Point of Failure

Q data item to be locked / unlocked

n number of replicas for a data item Q

(*) The availability discussion is limited to lock requests only. Even if a lock is granted, depending on the replication policy it may not be possible to write. For example, the majority protocol grants a write lock if more than half of the replicas are available, but the replication policy may require writes to be executed on all replicas, which requires *all* replicas to be available.