Department of Computer Science

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Similarity Search in Large Databases Wintersemester 2023/2024



Exam 31.01.2024

Name:	Student ID:	

Hints

- Check whether you received all pages of the exam (9 pages).
- Write your name or your student ID on each sheet of the exam and hand in all pages.
- All answers are expected to be written on the exam sheets.
- Clearly highlight and enumerate additional pages that are used for longer answers. Match your text with the according exercise.
- Only use pencils that are permanent and non-red colored.
- Use the notation and techniques discussed in the lecture.
- Exercises with more than one solution are not graded.
- You are allowed to use one A4 sheet with your personal notes (both sides, hand written or printed).
- Exam duration: 90 minutes

Signature

Grading

Filled by the examiner

Exercise	1	2	3	4	5	6	7	8	Summe
Total points	2	2	2	2	2	2	2	2	16
Points reached									

For a given distance function δ , the distance matrix d_{ij} for two sets of objects $A = \{a_1, a_2, a_3\}$ and $B = \{b_1, b_2, b_3\}$ is given in Figure 1.

Figure 1: Distance matrix d_{ij} for two sets of objects A and B.

	b_1	b_2	b_3
a_1	5	3	2
a_2	1	3	4
a_3	6	2	1

Compute the result tuples for the following matching strategies:

- a) Threshold Matching with $\tau = 2$.
- b) k-Nearest Neighbor Matching with k = 1.
- c) Global Greedy Matching.

Compute the edit distance between the strings lights and tight. Use the matrix produced by the dynamic programming algorithm to derive the shortest edit scripts and represent them with the gap representation. Consider the edit distance computation between two strings. One edit operation affects at most $q\ q\text{-}\mathrm{grams}.$

Provide an example where two edit operations affect strictly less than 2q q-grams.

2 Points

Exercise A Complexity	v of the	Trop Edit Distance Algorithm	2 Points
Exercise 4 - Complexity	y or the r	rree Eait Distance Algorithm.	2 Points

Prove that the number of key roots of a tree is equal to the number of leaves, i.e., |kr(T)| = |leaves(T)|.

Consider ordered trees T_1 and T_2 in Figure 2, forest distance matrix fd, and tree distance matrix td for the trees T_1 and T_2 .

- a) Compute the left-most leaf descendant arrays l_1 , l_2 for trees T_1 , T_2 , respectively.
- b) Fill the missing values for d_i and d_j into the forest distance matrix fd for the circled keyroot nodes in trees T_1 and T_2 .
- c) Circle the cell in the forest distance matrix that stores the distance between the prefixes $T_1[3..4]$ and $T_2[1..3]$.
- d) Cross all cells in the forest distance matrix and/or the tree distance matrix required to compute the distance between the prefixes $T_1[3..4]$ and $T_2[1..3]$.

fd	$\begin{vmatrix} \overrightarrow{d_j} \end{vmatrix}$			
$d_i \downarrow$				

td	1	2	3	4	5
1					
2					
3					
4					
5					
6					



Figure 2: Two ordered trees T_1 and T_2 .

For the ordered trees T_1 and T_2 in Figure 4:

- a) List the pq-gram profiles of T_1 and T_2 with p = 2 and q = 3.
- b) Compute the pq-gram distance of T_1 and T_2 .



Figure 3: Two ordered trees T_1 and T_2 .

2 Points

For the ordered trees T_1 and T_2 in Figure 4:

- a) Represent T_1 and T_2 as normalized binary trees and compute the binary branch distance.
- b) Based on the binary branch distance, what is the smallest value that the edit distance between T_1 and T_2 can adopt?



Figure 4: Two ordered trees T_1 and T_2 .

Consider the collection $R = \{s_1, s_2, s_3, s_4\}$ of sets in Figure 5. Build an inverted index by addressing the following tasks:

- a) Find a suitable ordering technique and sort the sets accordingly.
- b) Given a prefix of 2 for all sets, illustrate the inverted list index.

$$s_{1} = \{A, C, B, D, F, E\}$$

$$s_{2} = \{D, E, F, B, A\}$$

$$s_{3} = \{B, D, F, G\}$$

$$s_{4} = \{C, B, G\}$$

Figure 5: Set collection $R = \{s_1, s_2, s_3, s_4\}.$